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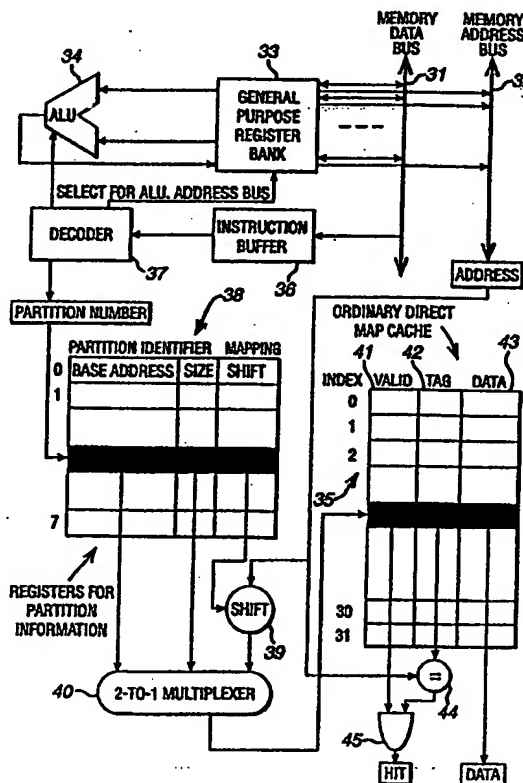
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(54) Title: CACHE MEMORY

(57) Abstract

A cache memory (35) has a logical organisation in which its memory space is divided into sub-sections or partitions (P). This permits different data objects to be allocated to different partitions during the operation of the cache memory (35). Commands used by the cache memory (10) may contain an extra parameter which is used to identify the appropriate partition within the cache memory (35). The parameter is extracted from the command by a decoder (37) and is passed to a specific line of an equator set (38) which contains identifiers which determine the partition to be used. To minimise data collisions for a given partition size, a stride may be defined which expresses the separation of addresses and from which a mapping function can be selected which covers all addresses in the cache memory (35) in an efficient way.



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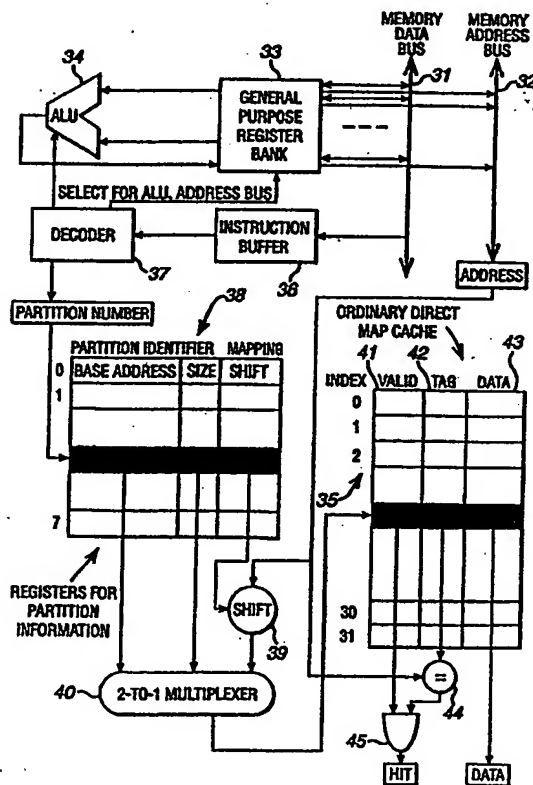
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A cache memory (35) has a logical organisation in which its memory space is divided into sub-sections or partitions (P). This permits different data objects to be allocated to different partitions during the operation of the cache memory (35). Commands used by the cache memory (10) may contain an extra parameter which is used to identify the appropriate partition within the cache memory (35). The parameter is extracted from the command by a decoder (37) and is passed to a specific line of an equator set (38) which contains identifiers which determine the partition to be used. To minimise data collisions for a given partition size, a stride may be defined which expresses the separation of addresses and from which a mapping function can be selected which covers all addresses in the cache memory (35) in an efficient way.



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CACHE MEMORYBACKGROUND OF THE INVENTIONFIELD OF THE INVENTION

5 The present invention relates to a cache memory and to a method of operating such a cache memory.

SUMMARY OF THE PRIOR ART

Current computer architectures rely heavily on the use of cache memory (hereinafter "cache"). Integrated with the processor on a single large chip, caches enable the processor to operate at high speed, as most instructions and data can be rapidly accessed from the caches instead of from the main memory which is usually at least ten times slower. On-chip caches have grown steadily in size over the last decade, and now represent a significant proportion of the cost and power consumption of the processor chip. It should be noted that the cache memory is inevitably of smaller memory space than the main memory, but provides more rapid access.

10
15
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Although it is normally the case that large caches offer better performance than small ones, it is also clear that the performance is not directly related to the size of the cache. Typically, program performance will increase as the cache size increases up to a certain

25

point at which further increases in cache size will have little or no effect. Cache management hardware takes no account of the characteristics of specific programs, and in many simple cases performs very inefficiently.

5 Another common problem is interference, which arises when a program accesses a collection of data objects which compete for parts of the cache. Current approaches to these problems have relied on the use of more complex cache architectures and on increasing cache sizes, with a
10 corresponding increase in system cost, size and power consumption.

SUMMARY OF THE INVENTION

At its most general, the present invention proposes
15 that a cache memory has a logical organisation in which its memory space is divided into sub-sections (hereinafter "partitions") under the control of a programmer or compiler. The size of the sub-sections need not be fixed, but may be determined by the control
20 operation.

This permits data objects to be allocated to particular partitions of the cache. This partitioning of the cache improves the performance of the cache such that a small cache memory can provide the same performance as
25 a conventional cache memory many times larger. This is useful because small caches are faster, take less chip

space and less power.

In addition, by minimising or eliminating interference, the performance of the cache and hence the program can be made more predictable.

5 In a first alternative, the partition to or from which data items are transferred is controlled by a parameter within an instruction such as a load or store instruction. The parameter may be different for different commands so that data items for different
10 commands made are of different partitions.

Thus, in a first aspect, the present invention may provide a method of operating a cache memory, using commands which cause a transfer, of corresponding items of data between the cache memory and a main memory, which
15 commands have an instruction component and an address component, the method comprising:

 - defining a plurality of sub-sections within the memory space of the cache memory, each of which has an associated identifier, the sizes of the sub-sections
20 being selectable from a range of sizes during the operation of the cache memory;

 extracting from the instruction component of a command a parameter corresponding to a selected one of the identifiers, the corresponding parameter being
25 different for different commands; and

 transferring items of data corresponding to said

command between the main memory and the sub-section of the memory space of the cache memory for which the associated identifier corresponds to the parameter of said command.

- 5 However, if registers are associated with instructions such as load or store, with a specific instruction being a corresponding register, then the parameter which determines the partition to or from which data items are transferred may be determined by such
- 10 registers themselves.

Thus, in a second aspect, the present invention may provide a method of operating a cache memory, using commands which causes a transfer, of corresponding items of data between the cache memory and a main memory, at

15 least some of which commands each have a corresponding register connected to a communication bus for use by said commands, the corresponding register being different for different commands, the method comprising:

defining a plurality of sub-sections within the

20 memory space of the cache memory, each of which has an associated identifier, the sizes of the sub-sections being selectable from a range of sizes during the operation of the cache memory;

associating a parameter with each said corresponding

25 register, each said parameter corresponding to a selected one of the identifiers; and

transferring items of data corresponding to said command between the main memory and the sub-section of the memory space of the cache memory for which the associated identifier corresponds to the parameter of the register corresponding to said command.

Another possibility for allocating data objects to particular partitions of the cache arises when a DMA controller is being used. Such a DMA controller generates specific commands for the memory access controlled by the DMA controller. Since the DMA controller generates those commands, it may also control the partition to or from which data items associated with those commands are transferred. Thus, in this case, the parameter which identifies the appropriate partition is not derived from the instruction, or a register associated with the instruction, but instead the command and its associated parameter are generated by a common trigger from the DMA controller.

Thus, in a third aspect, the present invention may provide a method of operating a cache memory under control of a DMA controller, the DMA controller being arranged to generate predetermined commands, the method comprising:

defining a plurality of sub-sections within the memory space of the cache memory, each of which has an associated identifier;

generating, at said DMA controller, one of said predetermined commands and a parameter associated with said one of said predetermined commands and with a selected one of the identifiers; and

- 5 transferring items of data corresponding to said one of said predetermined commands between a main memory and the sub-section of the memory space of the cache memory for which the associated identifier corresponds to the parameter of said one of said predetermined commands.
- 10 Preferably, the programmer or compiler is able to control the size of each partition. That permits analysis of the pattern of access to the cache, and division of the cache into suitably sized partitions, along with the derivation of an appropriate mapping
- 15 function to map memory addresses to addresses of lines in the partition. Once that has happened, the mapping function should be able to map items in a data structure which are accessed in sequence onto different lines within a partition which the compiler uses for that
- 20 structure. The aim is then to minimise data collisions for a given partition size. To do this, it is preferable to derive from the program a quantity hereafter referred to a "stride", the value of which defines the separation of the addresses within the address space of the main
- 25 memory of successive accesses to or from the memory.
- Based on the stride, a mapping function can be selected

that generates addresses which cover all addresses within a cache partition in an efficient way.

Thus, in a fourth aspect, the present invention may provide a method of operating a cache memory, comprising:

5 defining a plurality of sub-sections within the memory space of the cache memory; and

transferring data items associated with each other only to a corresponding sub-section of the memory space of the cache memory;

10 wherein each sub-section has a stride associated therewith, the stride representing the separation within the memory addresses of a main memory of successive transfers of data between the corresponding sub-section and the main memory.

15 In each of the above four aspects, the present invention may also provide a memory system arranged to operate as discussed above.

It should be noted that although such control of the partitioning is preferable, a general purpose function
20 may be needed to perform mapping, if e.g the pattern of access to the data is not known to the compiler.

Preferably, the compiler controls the partitioning of the cache memory using a parameter added to a load and store instruction. That partition parameter may be derived
25 from the instruction opcode or from one or more registers. Each of these registers may be a general

purpose register, or may be one or more dedicated partition registers. Registers that are used to implicitly access memory, eg via stack pointer or program counter, normally will have a dedicated partition register associated with them.

It is usually desirable that there are functions which identify the line of cache memory from the memory address, and in this case it is preferable that each partition has its own function. The function may for example be a shift and modulo operation.

As has been mentioned above, the stride defines the separation of successive accesses to the memory, for each partition. It should be noted that multiple partitions may be used to cache accesses with different strides to the same data object.

With the present invention it is possible for multiple DMA controllers and a processor to use a common cache, by providing a dedicated partition register in each controller so that the different controllers and the processor all access different partitions.

BRIEF DESCRIPTION OF THE DRAWINGS

An embodiment of the present invention will now be described in detail, by way of example, with reference to the accompanying drawings, in which:

Figs. 1a to 1c show a cache memory according to the

present invention, which is dividable into partitions:

Figs. 2a and 2b show the structure of the partition information;

Fig. 3 is a schematic block diagram one arrangement
5 of a memory structure usable in the present invention;

Fig. 4 is a table showing partitions which may be used by a specific program in an embodiment of the present invention;

Figs. 5a to 5d are computer code fragments for
10 controlling a partitioned cache;

Fig. 6 is a schematic block diagram of another arrangement of a memory structure usable in the present invention;

Fig. 7 is a schematic block diagram of yet another
15 arrangement of a memory structure usable in the present invention; and

Fig. 8 is a schematic block diagram of a further arrangement of a memory structure usable in the present invention.

20

DETAILED DESCRIPTION

Consider a cache memory (cache) with 2^n lines, each line being able to store items in a data structure. Fig. 1a shows such a cache 10 where n is 5, so that there are
25 32 lines. Such a cache 10 can be divided into 2^p partitions each of size $2^{(n-p)}$ lines. It can also be

divided into partitions of size $2^{(n-x)}$ where $x < p$, or into any combination of different size partitions. Each partition P has an address P_a which corresponds to the address of the first line in the partition, so that an
5 address a within partition P has an address formed by carrying out a bitwise OR operation on P_a and a . The address of the line to be used within the partition must, of course, be derived from the memory address used to load or store data. Thus, Fig. 1b shows the cache of
10 Fig. 1a divided into 8 partitions 10a to 10h of equal size, and Fig. 1c shows the same cache 10 divided into unequal partitions 10j to 10n.

In such an arrangement, a program produced by a compiler or programmer must be able pass information to
15 the hardware of the cache. In an embodiment of the present invention, it is proposed that an extra parameter be added to the load and store instructions which control the cache memory. This extra partition parameter supplies the partition information for operation of the
20 cache. This will be discussed in more detail subsequently.

However, it should be noted that there are two alternatives which do not require the instruction set to be modified. It is possible use the high order bits of
25 the address space to contain the partition information, or alternatively it is possible to store the "current

partition information" in the cache, and change this partition information as and when required. The last solution is only of use when it is expected that a series of multiple requests will be sent through the same
5 partition. Using higher order address bits is particularly useful when a cache according to the present invention is to be used together with an existing processor core. It is also suitable for languages like C, where the partition information will be carried along
10 implicitly with a pointer.

The structure of the partition information also needs to be considered. A simple method is to use a number as the partition operand and to use this number to select one of a set of partition control registers
15 holding partition identifiers and mapping information. The partition identifier can be represented by the partition address and partition size. A more complicated but more elegant scheme is to use a bit-pattern in which the position of the rightmost "one" bit defines the size
20 of the partition and the leftmost bits define the address of the partition within the cache. This is depicted in Figs. 2a and 2b, in which Fig. 2a shows a general case of division into a partition address 20 and a partition size 21, and Fig. 2b shows a specific case. The mapping
25 information defines how to hash the address in the partition. In our simple scheme, this information

consists of the shift length, but a more complicated scheme might require an XOR of some parts of the address.

One possible scheme currently preferred is to use numbered partitions, and pass an extra parameter with 5 instructions accessing the memory. Assuming a normal RISC load/store architecture, the instructions affected are the load and store instructions. For our example we have LOAD and STORE that offer indirect load and store operations. Each of these instructions requires an extra 10 parameter (compared with a normal RISC load/store), which indicates the partition number in the cache. It would of course be possible to include other instructions such as indexed loads and stores.

We create partitions using a CPART instruction. The 15 CPART instruction takes 3 parameters, a partition number, a number of lines and a stride. The partition number is the name that will be used in the load and store instructions. The number of lines (currently we restrict this to powers of two) is the size of the partition, and 20 the shift indicates the shift which should be applied to the memory address before it is used to access this partition.

The implementation of such a partitioned cache can be based on a conventional direct map cache memory. A 25 block diagram of one arrangement of a suitable structure is shown in Fig. 3.

In a conventional memory structure, a command to transfer data between a memory and, for example, a register for subsequent use includes an address component, and an opcode. Examples of such commands are

5 load commands which retrieve data from a specified address in a memory, and store commands which store data at a specified address in a memory. As shown in Fig. 3, the address and data are normally carried by two separate busses 31, 32, one 31 of which carries data including an

10 instruction component of the command and the other 32 of which carries the address. Signals to and from those buses 31, 32 may pass via a register bank 33, and from there to an arithmetic logic unit (ALU) 34.

In the present invention, however, the instructions

15 on bus 31 contain an extra parameter which is used to identify the appropriate partition within a cache memory 35. Such an instruction is then passed via a buffer 36 to a decoder 37 which extracts from the instruction the parameter, which is output from the decoder 37 as a

20 partition number. That partition number is then passed to a register set 38 which contains partition information.

As shown in Fig. 3, that register set 38 contains identifiers as described with reference to Figs. 2A and

25 2B, identifying a base address, a size and a shift. The partition number output from the decoder 37 identifies a

specific line within the register set 38, to output a partition base address, size and shift. The latter is used to perform a shift on the address derived from the address buss 32 to a shifter 39, the output of which is
5 fed to a multiplexer 40.

The multiplexer also receives the partition identifier from the register set 38, being the base address and size of the partition identified by the partition number from decoder 37. The output of the
10 multiplexer then identifies the appropriate line of the cache 35. The cache 35 is a direct map cache, with each line divided into a validity bit or bits 41, tag bits 42 and data bits 43. When the signal from the multiplexer 40 identifies a line in the cache 35, the tag 42 is
15 output to a comparator 44, which compares the tag with the address from the address bus 32. If equality is found, an output is sent to an AND gate 45, which also received an input from the validity bit 41. The logical AND operation then confirms that the appropriate line has
20 been identified, and the data 43 can then be read.

In such an embodiment, the addressing of the partition control registers can be pipelined with the execution of the load/store instruction if we assume that the partition operand is a constant parameter of the load
25 and store instructions. Also, the partition control registers can simply be general purpose registers.

The partition control register set can be very small indeed. Normally one register is needed for each partition. The base address and select bits can be stored in $\log_2 l$ bits where l is the number of lines in the cache. These may be combined in one word of size $1 + \log_2 l$ bits if the encoding presented earlier is to be used. Finally, the shift needs to be stored in at most 6 bits for 64 bit address machines.

In this embodiment, suppose that the compiler takes a program which uses scalars and (multi dimensional) arrays, and generates instructions which include all the partitioning information. The compiler may then calculate the minimally required partition sizes, and analyse all accesses in order to optimise persistence (the length of time each value remains) in each of the partitions.

The "stride" of a reference is the distance between addresses of successive accesses to an array variable.

For example,

```
20      int i, j, k;
        int cTemp, dTemp;
        int [ 32, 32 ] a;
        int [ 32, 32 ] b;
        int [ 32, 32 ] c;
25      int [ 32, 32 ] d;
```

is a set of variable declarations for variables a, b, i,

16

```

j and k, and
    for i = 0 to 31 {
        for j = 0 to 31 {
            cTemp = 0;
5           dTemp = 0;
            for k = 0 to 31 {
                cTemp = cTemp + (a[k,i] * b[i,k]);
                dTemp = dTemp + (b[k,i] * a[i,k]);
            }
10          c[i,j] = cTemp;
            d[i,j] = dTemp;
        }
    }

```

is a sequence of statements, which calculates AB and
 15 BA for two matrices A and B in matrices C and D. All
 matrices are stored in two dimensional arrays of size 32
 by 32, which are stored as sequences of 1024 values. The
 first value denotes element a[0,0], the second value
 element a[0,1], ..., the 32-nd value is element a[0,31],
 20 the 33-rd value is element a[1,0] and so on until the
 1024th value which is a[31,31]. Note that a[0,0] and
 a[0,1] are one memory cell apart, while a[0,0] and a[1,0]
 are 32 memory cells apart.

In general, we can handle any indexed variable where
 25 the index is of the form $c_0k + c_1$. Here c_0 and c_1 are
 constants and k is the loop counter. In the line

dtemp = dtemp + b [k,i] * a [i,k]

a has stride 1 and b has stride 32.

Suppose that an array is to be accessed repeatedly within a loop with a stride s , by which it is meant that
5 successive accesses are to elements s apart. By extracting bits from the addresses used, starting at the bit position defined by the least significant 1-bit of s , line addresses can be generated which will change with each iteration of the loop, distributing the data
10 accessed by successive iterations among the lines within the cache partition.

To see that this will in fact distribute the data optimally, consider a stride s . If we ignore all trailing zeros of s are ignored, a stride s' is obtained
15 which must be odd. The stride s' is either 1 or is co-prime with any power of 2. Therefore, $k \times s' \bmod 2^c$ ($k=0,1,2,\dots;c$) will traverse all numbers between 0 and 2^c-1 . Hence all strides s use every line in the partition before reusing any of them. The partition size will
20 therefore define the persistence of data within the cache.

For example, if the array which is accessed with a stride of 20, (binary 10100₂) then a shift down by two bits may be used to map array addresses to line addresses
25 within the partition used for the array accesses. If the size of the partition is 8, then subsequent items will be

placed at line addresses 0,5,2,7,4,1,6,3,0,5,..... The persistence of data is 8 in this case; if this is not enough, it may be necessary to increase the partition size.

5 In this embodiment, the compiler may create one special partition for the scalars, and then one partition for every group of accesses to an array which have identical strides. The scalar partition assumes that scalars are placed contiguously in memory, and start at a
10 cache line boundary. The number of lines allocated to the scalar partition is determined by the number of scalars. By default a partition is allocated which is large enough to hold all the scalars. In the example above there are 5 scalars. The cache has 4 words per
15 line. Hence, a two line partition is allocated for the scalars with the instruction

CPART 0, 2, 1

All scalar references are then marked to use partition 0 when loading and storing.

20 For all groups of non-scalar references, a unique partition is created. The size of this partition depends on the mapping function that the cache uses to access elements in the partition and on the required persistence.

25 The mapping function used selects the line in the partition by taking all bits from the access address

starting from the position of the least significant 1 bit in the stride (as defined previously). The compiler allocates partition sizes large enough to keep each data item in the partition until it is no longer needed.

- 5 In the example above, the complete partition summary is contained in the table in Fig. 4. Note that the total cache size allocated for this is 8 cache lines, spread over 7 partitions.

Various possibilities for controlling a cache memory
10 according the present invention by software are as illustrated in Figs. 5a to 5c. In those arrangements, the code fragments will perform a vector addition, with the vectors being stored with different strides for illustrative purposes.

- 15 In the first code fragment of Fig. 5a, the semantics of the instructions are as follows:

CPART p,b,l,s This instruction creates partition number
 p. The partition starts at line b in the cache, and
 consists of l lines. The last parameter indicates
20 the shift value, the base address will be shifted by
 s before indexing the direct mapped cache.

- LOAD d,[s],p This instruction loads the contents of the
 memory address pointed to by s into d, where s and d
 are general purpose registers. In this example we
25 have restricted ourselves to indirect loads via a
 register, but this could be any addressing mode.

The third parameter is the partition number via which the load should be performed. The partition should be defined with the CPART instructions before using it in a load.

5 STORE *d*, [*s*], *p* This instruction stores *d* in the memory address pointed to by *s*, where *s* and *d* are general purpose registers. In this example we have restricted ourselves to indirect loads via a register, but this could be any addressing mode.

10 The third parameter is the partition number via which the load should be performed. The partition should be defined with the CPART instruction before using it in a load

ADD, SUB, BGT perform addition, subtraction and a

15 conditional branch, similar to conventional processors. Only load and store instructions need an extra parameter.

In the second code fragment shown in Fig. 5b the
20 CPART instruction has a different semantics:

CPART *p*, *i* This instruction creates partition number *p*.

The base address, size and shift are all mapped into a single number. The last six bits of this number
25 are the shift (a number between 0 and 63), the bits before that are encoded using the scheme discussed

in figure 2.

The code fragments described with reference to Figs. 5A and 5B are used with the memory structure of Fig. 3, in which there is a register set 38 which stores cache partition information. Fig. 6 illustrates a modification of the memory structure of Fig. 3. In Fig. 6, components which correspond to the components in Fig. 3 are indicated by the same reference numbers. In Fig. 6, however, the partition information derived by decoder 37 from the instructions received from instruction buffer 36 are passed to the register bank 33 in which the partition and shift information is stored. Storage in the register bank 33 may otherwise be the same as in Fig. 3. Thus, the register bank 33 outputs partition information 50 containing a base address, size and shift which are passed to the multiplexer 40 and shifter 39 respectively. The arrangement of Fig. 6 is then otherwise the same as that of Fig. 3. However, the commands needed are then changed, and the code fragment for this is illustrated in Fig. 5C.

Thus, in a third code fragment shown in Fig. 5C conventional registers are used instead of partition registers. Three conventional registers are loaded with the partition address, size, and shift, and these registers are used as the partition operands of the load

and store operations. The third parameter of the load and store is now a register parameter:

LOAD d,[s],p the contents of the partition register
5 associated with the register S determines the base address, size and the shift information for the partition.

STORE d,[s],p the contents of the partition register associated with the register S determine the base
10 address, size and the shift information for the partition.

Fig. 7 shows another alternative, in which partition registers 60 are associated directly with the register
15 bank 33 contain the partition and shift information. Again, the components in Fig. 7 which are the same as those in Figs. 3 and 6 are indicated by the same reference numerals. The code fragment in such an arrangement is illustrated in Fig. 5d.

20

LOAD d,[s],p p is a register, the contents of which determine the base address, size and the shift information for the partition.

STORE d,[s],p p is a register, the contents of which
25 determine the base address, size and the shift information for the partition.

Fig. 8 shows a further alternative, which is applicable when a DMA controller 70 is to access the cache 35. The rest of the structure of Fig. 8 is the same as that in Fig. 3, and corresponding parts are indicated by the same reference numerals.

When the DMA controller 70 operates, it generates commands which are transmitted via the busses 31, 32. As shown in Fig. 8, the DMA controller 70 has a control logic unit 71 which signals to data registers 72, address registers 73 and counters 74 to generate such commands. The data for the commands from the data register 72 are passed to the bus 31, and the corresponding addresses is from the address register 73 to the bus 32. Since the DMA controller 70 generates such commands, it is possible for the DMA controller 70 directly to determine which partition of the cache 35 needs to receive or output data for each command generated by the DMA controller 70. Therefore, the control logic unit 71 may, at the same time as it triggers a command, may cause a partition number unit 75 to output data representing a partition number, which is sent to the register set 38. Once the partition number has thus been supplied, access to or from the cache 35 occurs in the same as in Fig. 3.

It can be noted that, in such an arrangement, it is possible that only some of the partitions of the cache 35 are operated under the controller of the DMA controller

70. Then, other partitions may be accessed from commands on the busses 31, 32, via the instruction buffer 36 and the decoder 37, as in the arrangement of Fig. 3.

Investigations have shown that a partitioned cache
5 in accordance with the present invention can achieve results comparable with larger caches of known configuration. Because a cache according to the present invention is physically smaller, the cost of production is reduced. Also, there may be more physical space
10 available for other associated electronic components in spatially restricted devices. A further advantage is that, because a cache according the present invention may be smaller than an equivalent conventional cache, it may consume less power.

15 It is possible to implement a partitioned cache according to the present invention in several ways. The embodiment of Fig. 3 uses a direct mapped cache, which conventionally is built from memory elements that are faster and more expensive than main memory. Because, for
20 equivalent functionality, a cache according to the present invention may be smaller, it may be implemented using even faster, register style memory devices. This means that even when the overall hit ratio of a cache according to the present invention is lower than known
25 caches, the disadvantage is not significant because of the improvement in speed of access to the cache which is

gained by using register style memory.

Because the present invention permits various activities to be kept in different partitions of the cache, the performance of a partition is independent of
5 other partitions. This may be important for e.g audio and video applications. Moreover, predictability of performance is improved, because performance of the system is directly related to the performance of its constituent parts, namely the partitions. In the
10 conventional cache memory, an attempt to combine two functions in a programme may lead to faulty results.

In a conventional cache, data items dynamically compete for space. Using a cache partition according to the present invention, a compiler can allocate data items
15 to independent partitions, and thus the compiler has control over cache allocation in same way it has control over register allocation.

CLAIMS

1. A method of operating a cache memory, using commands which cause a transfer, of corresponding items of data
5 between the cache memory and a main memory, which commands have an instruction component and an address component, the method comprising:

defining a plurality of sub-sections within the memory space of the cache memory, each of which has an
10 associated identifier, the sizes of the sub-sections being selectable from a range of sizes during the operation of the cache memory;

extracting from the instruction component of a command a parameter corresponding to a selected one of
15 the identifiers, the corresponding parameter being different for different commands; and

transferring items of data corresponding to said command between the main memory and the sub-section of the memory space of the cache memory for which the
20 associated identifier corresponds to the parameter of said command.

2. A method according to claim 1, wherein each sub-section of the memory space of the cache memory has a predetermined rule defining in which line of the
25 corresponding sub-section successive items of data are stored, which predetermined rule makes use of the address

component of said command.

3. A method according to claim 2, wherein each predetermined rule is a shift and select operation.

4. A method according to claim 1, wherein each sub-
5 section has a stride associated therewith, the stride representing the separation within the memory addresses of the main memory of successive transfers of data between the corresponding sub-section and the main memory.

10 5. A method according to claim 4, wherein the lowest 1-bit in said stride determines the shift applied to the memory addresses of the main memory to determine the location within the corresponding sub-sections.

6. A method according to any one of the preceding
15 claims, wherein each parameter identifies one of a set of registers, each of which contains data which identifies the sub-section for which the associated identifier corresponds to that parameter.

7. A method according to claim 6, wherein the address
20 and/or the parameter of said commands are stored in general purpose registers.

8. A method according to claim 7, wherein the set of registers and the general purpose registers form a common unit.

25 9. A method according to any one of claims 1 to 5, wherein each parameter represents data which identifies

the sub-section for which the associated identifier corresponds to that parameter.

10. A method of operating a cache memory, using commands which cause a transfer of corresponding items of data
5 between the cache memory and a main memory, at least some of which commands each have a corresponding register connected to a communication bus for use by said commands, the corresponding register being different for different commands, the method comprising:

10 defining a plurality of sub-sections within the memory space of the cache memory, each of which has an associated identifier, the sizes of the sub-sections being selectable from a range of sizes during the operation of the cache memory;

15 associating a parameter with each said corresponding register, each said parameter corresponding to a selected one of the identifiers; and

transferring items of data corresponding to said command between the main memory and the sub-section of
20 the memory space of the cache memory for which the associated identifier corresponds to the parameter of the register corresponding to said command.

11. A method of operating a cache memory under control of a DMA controller, the DMA controller being arranged to
25 generate predetermined commands, the method comprising:

defining a plurality of sub-sections within the

memory space of the cache memory, each of which has an associated identifier;

generating, at said DMA controller, one of said predetermined commands and a parameter associated with
5 said one of said predetermined commands and with a selected one of the identifiers; and

transferring items of data corresponding to said one of said predetermined commands between a main memory and the sub-section of the memory space of the cache memory
10 for which the associated identifier corresponds to the parameter of said one of said predetermined commands.

12. A method of operating a cache memory, comprising:

defining a plurality of sub-sections within the memory space of the cache memory; and

15 transferring data items associated with each other only to a corresponding sub-section of the memory space of the cache memory;

wherein each sub-section has a stride associated therewith, the stride representing the separation within
20 the memory addresses of a main memory of successive transfers of data between the corresponding sub-section and the main memory.

13. A memory system comprising a main memory and a cache memory, the memory space of the cache memory being
25 divided into a plurality of sub-sections, each sub-section having a corresponding identifier, the sizes of

the sub-sections being selectable from a range of sizes during the operation of the cache memory;

instruction and address buses respectively for instruction components and address components of commands
5 which cause a transfer of data between the cache memory and the main memory;

means for extracting from the instruction component of a command a parameter corresponding to a selected one of the identifiers, the corresponding parameter being
10 different for different commands; and

means for controlling the transfer of data to and from the cache memory such that data corresponding to a command is transferred between the main memory and the sub-section of the memory space of the cache memory for
15 which the associated identifier corresponds to the parameter of said command.

14. A memory system comprising a main memory and a cache memory, the memory space of the cache memory being divided into a plurality of sub-sections, each sub-
20 section having a corresponding identifier, the sizes of the sub-sections being selectable from a range of sizes during the operation of the cache memory;

instruction and address buses respectively for instruction components and address components of commands
25 which cause a transfer of data between the cache memory and the main memory;

registers connected to said address bus, each of said registers being for a corresponding command, the corresponding register being different for different commands;

- 5 a parameter associated with each said register, each parameter corresponding to a selected one of the identifiers; and

means for controlling the transfer of data to and from the cache memory such that data corresponding to a
10 command is transferred between the main memory and the sub-section of the memory space of the cache memory for which the associated identifier corresponds to the parameter of the register corresponding to said command.

15. A memory system comprising:

- 15 a main memory;

a cache memory, the memory space of which is divided into a plurality of sub-sections, each sub-section having a corresponding identifier;

- a DMA controller having means for generating
20 predetermined commands, which commands involve transfer of data to or from the cache memory, and means for generating parameters corresponding to at least one of said identifiers, said parameters also being associated with said commands such that each command has an
25 associated parameter; and

means for controlling the transfer of data between

32.

the cache memory and the main memory such that data corresponding to a command is transferred to or from the sub-section of the memory space of the cache memory for which the associated identifier corresponds to the

5 parameter of said command.

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FIG. 1a

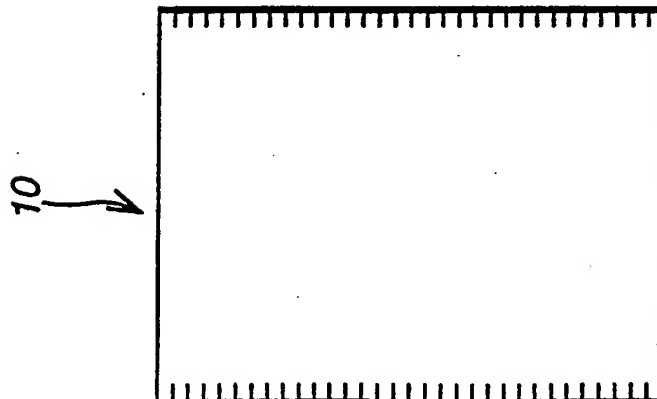


FIG. 1b

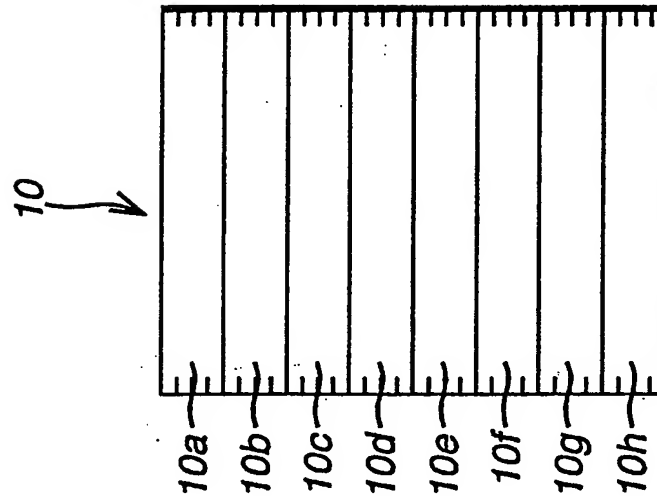


FIG. 1c

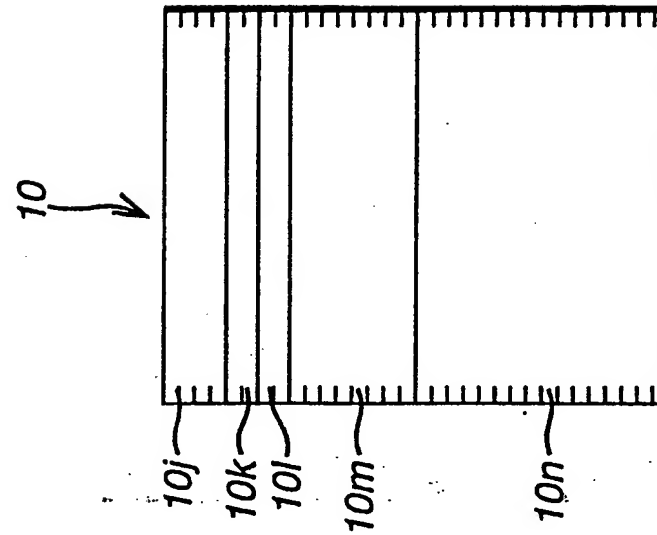
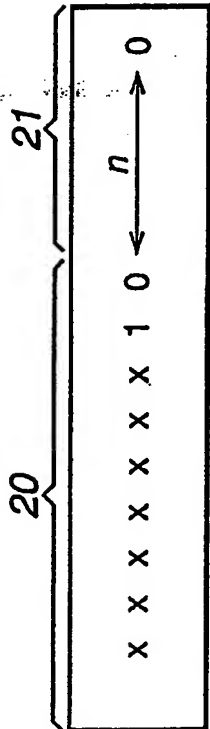
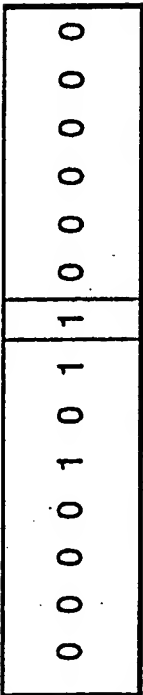


FIG. 2a



Partition Address: xxxxxx 2^n Partition Size: 2^n

FIG. 2b

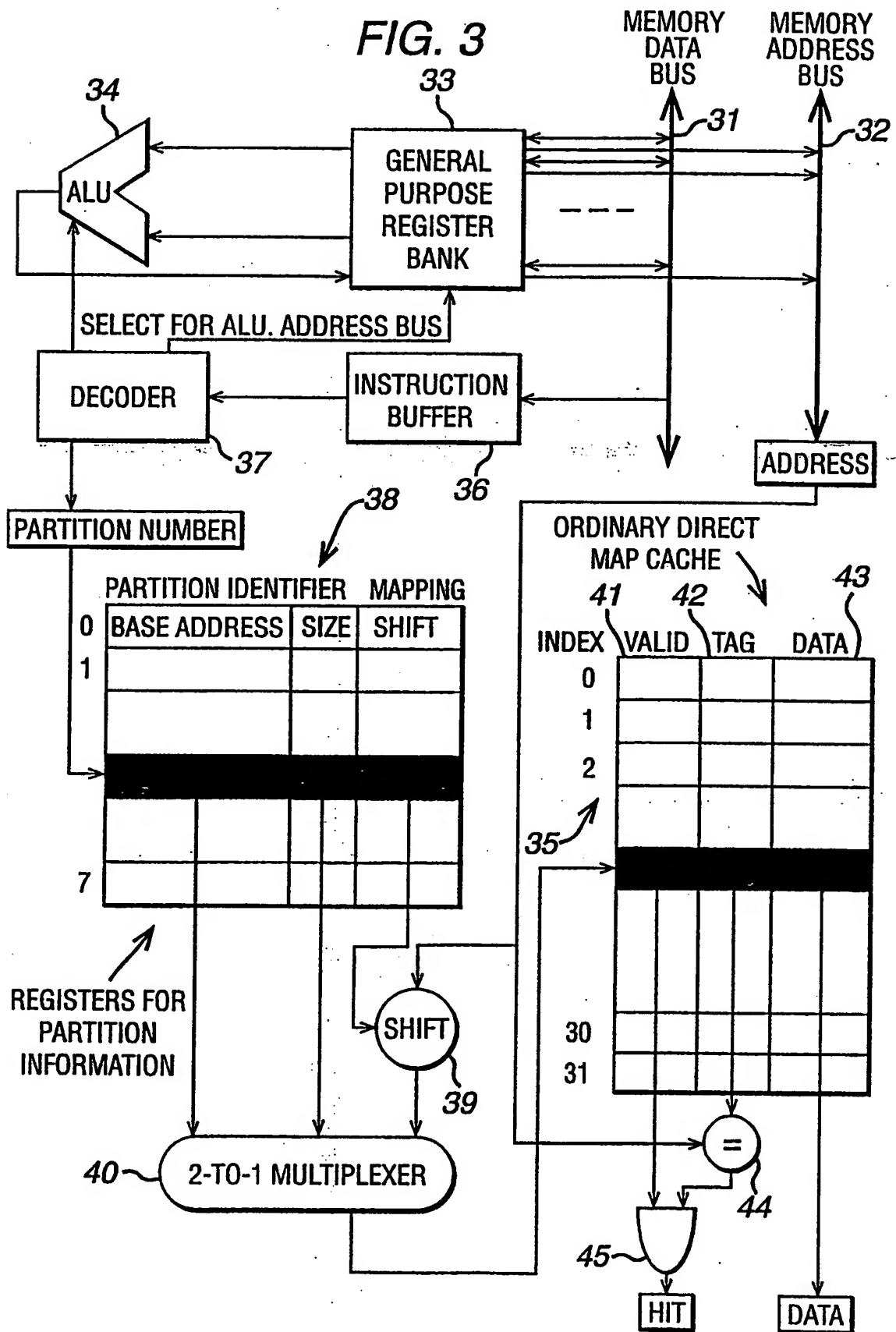


Partition Address: 320 Partition Size: 64

FIG. 4

Partition	Size	Stride	Contains
0	2	1	Scalar partition
1	1	1	Partition for d
2	1	1	Partition for c
3	1	1	Partition for b
4	1	32	Partition for b
5	1	32	Partition for a
6	1	1	Partition for a

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FIG. 3



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FIG. 5b

CPART	3, 0xA81
CPART	4, 0xC83
CPART	5, 0x741

LOOP:

LOAD	r4, [r0], 5
LOAD	r5, [r1], 4
ADD	r5, r5, r4
STORE	r5, [r2], 3
ADD	r0, r0, #4
ADD	r1, r1, #48
ADD	r2, r2, #8
SUB	r6, r6, #1
BGT	LOOP

FIG. 5a

CPART	3, 10, 2, 1
CPART	4, 12, 2, 3
CPART	5, 14, 1, 1

LOOP:

LOAD	r4, [r0], 5
LOAD	r5, [r1], 4
ADD	r5, r5, r4
STORE	r5, [r2], 3
ADD	r0, r0, #4
ADD	r1, r1, #48
ADD	r2, r2, #8
SUB	r6, r6, #1
BGT	LOOP

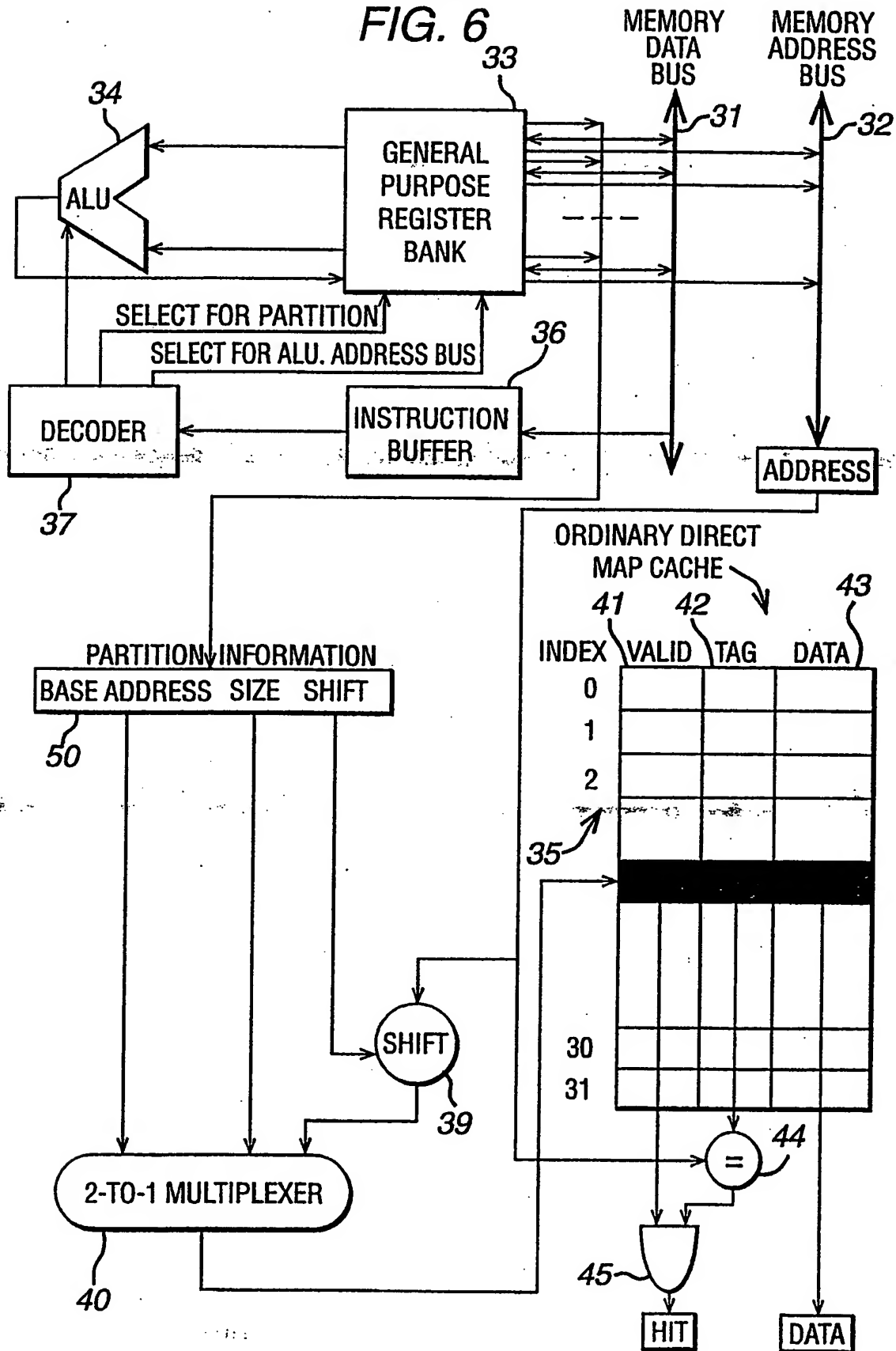
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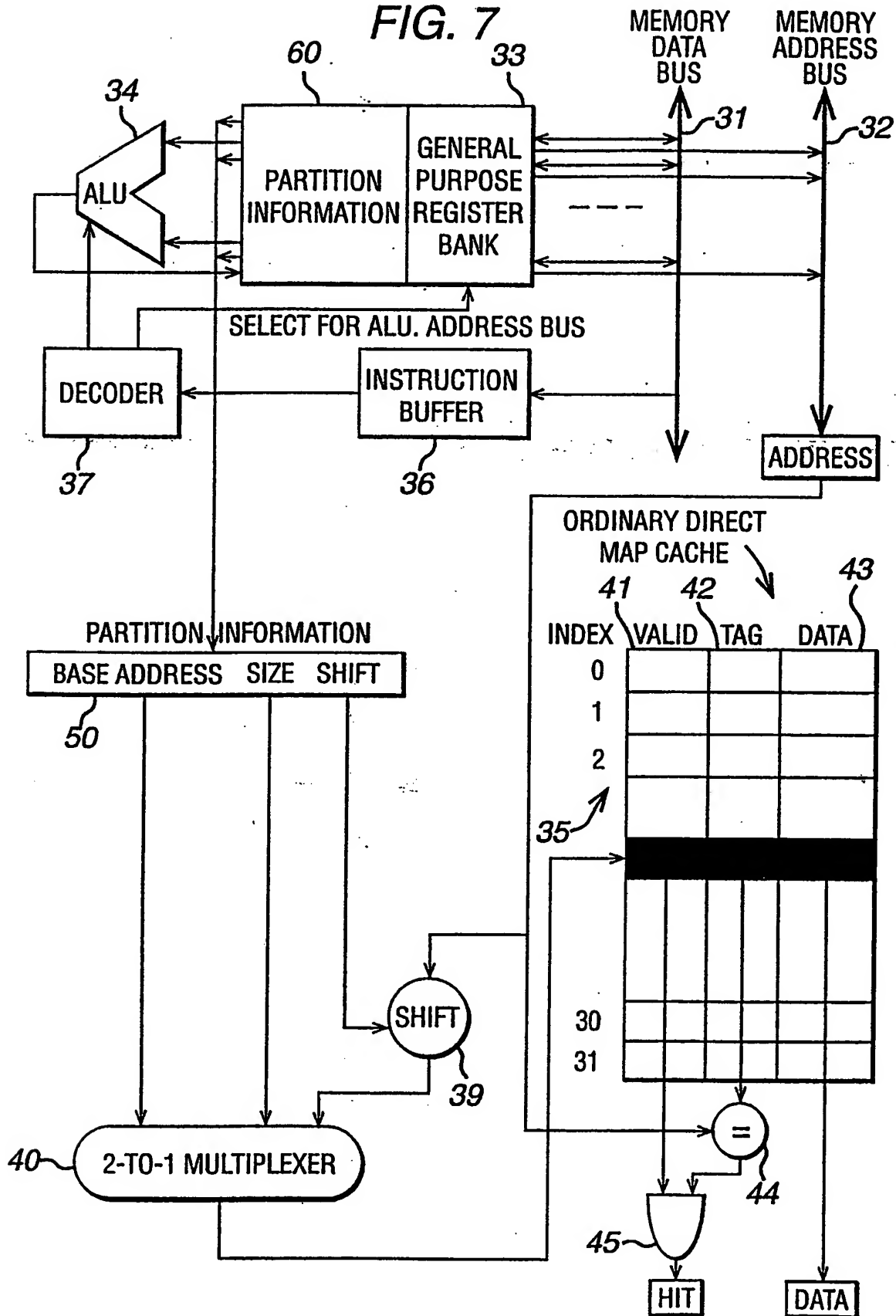
FIG. 5d

CPART	r2 0xA81
CPART	r1, 0xC83
CPART	r0, 0x741
LOOP:	
LOAD	r4, [r0], r10
LOAD	r5, [r1], r9
ADD	r5, r5, r4
STORE	r5, [r2], r8
ADD	r0, r0, #4
ADD	r1, r1, #48
ADD	r2, r2, #8
SUB	r6, r6, #1
BGT	LOOP

FIG. 5c

LDC	r8, 0xA81
LDC	r9, 0xC83
LDC	r10, 0x741
LOOP:	
LOAD	r4, [r0], r10
LOAD	r5, [r1], r9
ADD	r5, r5, r4
STORE	r5, [r2], r8
ADD	r0, r0, #4
ADD	r1, r1, #48
ADD	r2, r2, #8
SUB	r6, r6, #1
BGT	LOOP

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FIG. 6

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FIG. 7

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FIG. 8

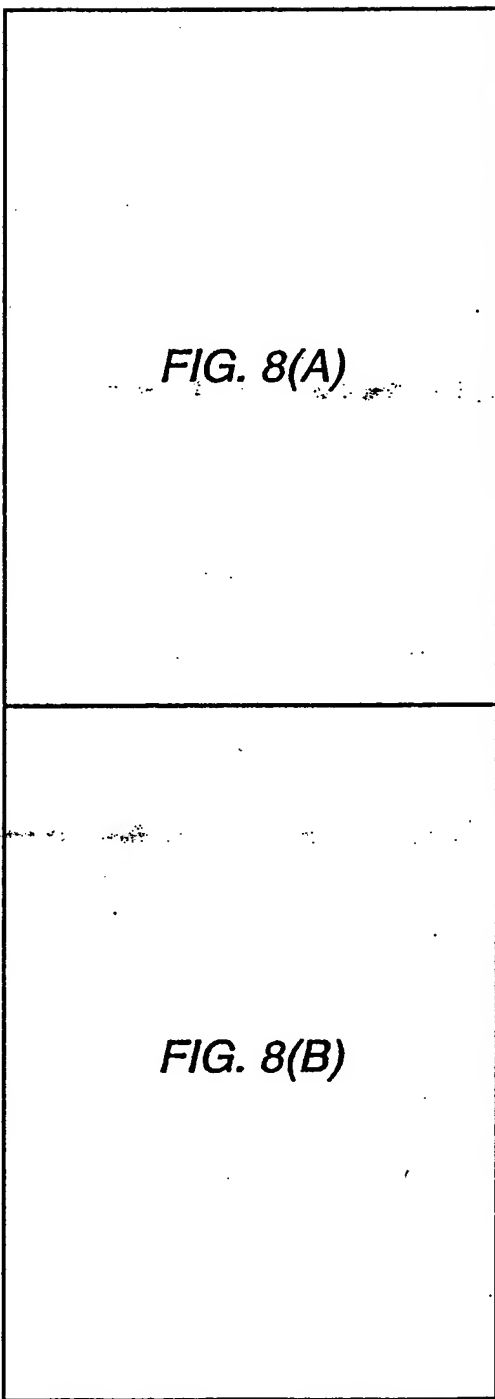
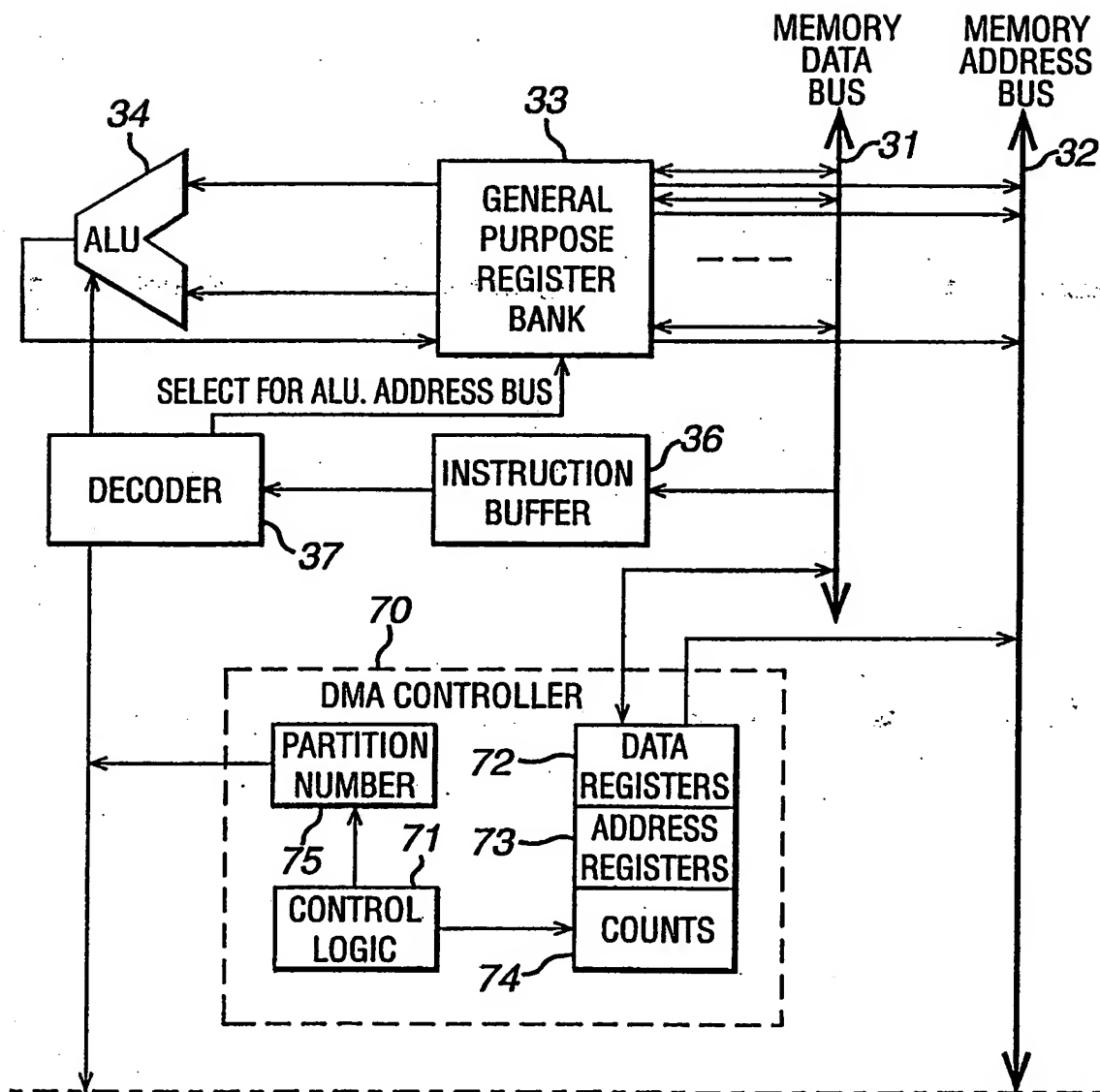


FIG. 8(A)

FIG. 8(B)

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FIG. 8(A)



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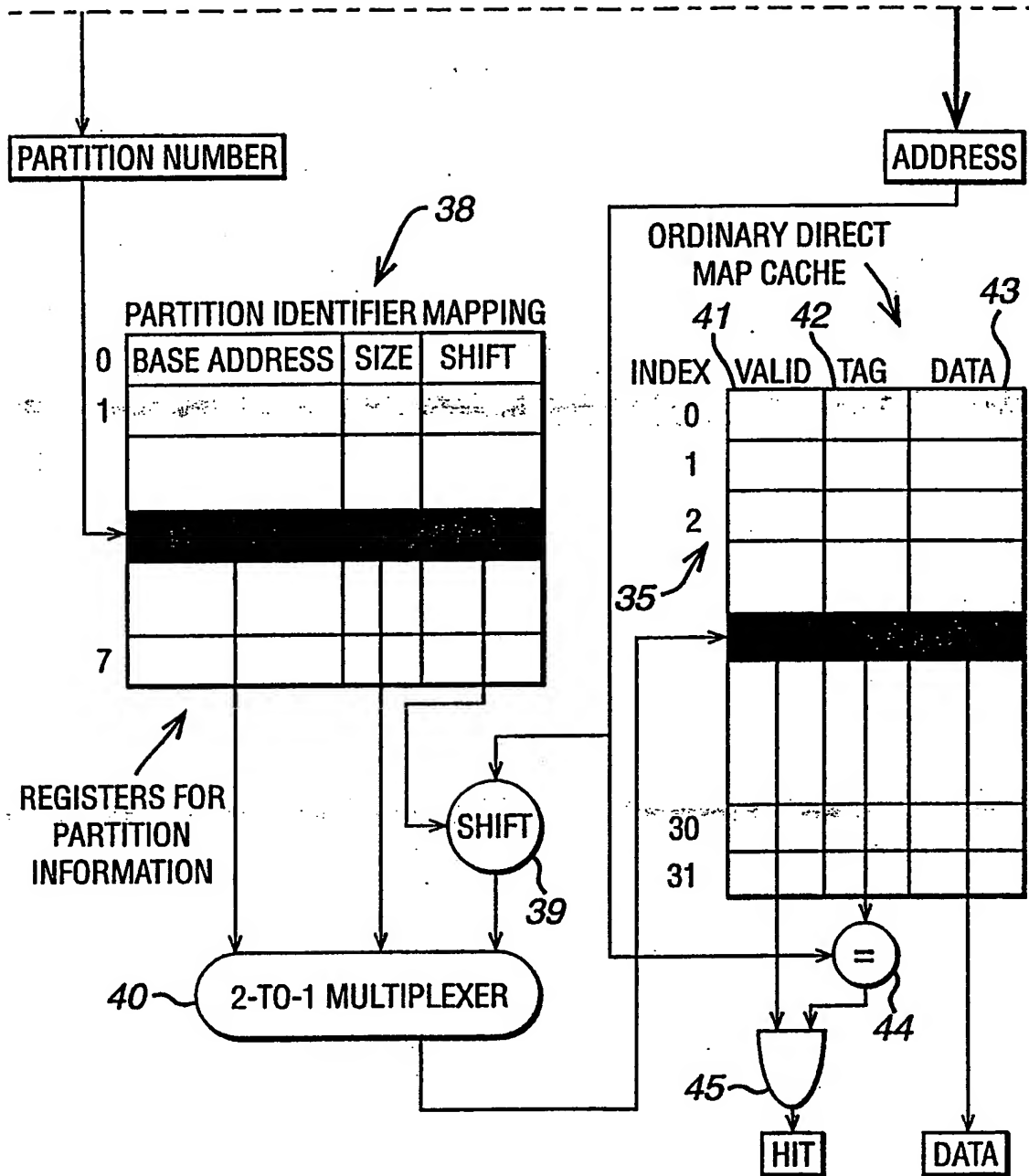


FIG. 8(B)

INTERNATIONAL SEARCH REPORT

International Application No

PCT/GB 00/00250

C.(Continuation) DOCUMENTS CONSIDERED TO BE RELEVANT

Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
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Information on patent family members

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